

## **Distinguishability of Random Quantum States**

Based on *Montanaro, A. Commun.Math. Phys. 3, 619–636 (2007).*

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## <span id="page-1-0"></span>**[Introduction](#page-1-0)**

<u> 1989 - Johann Barnett, mars et al.</u>

Given an unknown state  $\rho_? \in \mathcal{L}(\mathbb{C}^d)$ , picked from a known set of states  $\mathcal{E} = \{\rho_1, \ldots, \rho_n\}$  with a known prior probability distribution on  $\mathcal{E}$ ,

We want to find an **optimal** measurement to determine *ρ*?,

In the sense that the probability of success is optimized.

We can focus on finding a POVM measurement (Why?).

Given an oracle implementing an unknown *n*-bit Boolean function *f* :  $\{0,1\}$ <sup>*n*</sup>  $\mapsto$   $\{0,1\}$  picked uniformly at random from a known set *F* of functions,

Identify *f* with the minimum number of calls to the oracle.

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$$
|\psi_f\rangle=\frac{1}{2^{n-1}}\sum_{x=0}^{2^n-1}(-1)^{f(x)}|x\rangle
$$

## **Some Notations**

For  $\rho_i \in \mathcal{E}$ , which appears with probability  $p_i$ , define

$$
\rho_i':=p_i\rho_i.
$$

**If**  $ρ<sub>i</sub> = |ψ<sub>i</sub>⟩$   $\langle ψ<sub>i</sub>|$ , we define

$$
|\psi'_i\rangle = \sqrt{p_i} |\psi_i\rangle.
$$

For a measurement  $M = \{M_i\}_i$ , we denote the probability of success in distinguishing which state is given, by  $P^{\text{M}}(\mathcal{E}).$ 

$$
P^{opt}(\mathcal{E}) := \sup_{M} P^{M}(\mathcal{E})
$$

$$
P^M(\mathcal{E}) = \sum_i \text{tr}(M_i \rho_i')
$$

What are the discrimination strategies?

- non-measurement strategy: The probability of sucess is  $\sum_i p_i^2$ .
- the most natural way to design a measurement:

$$
M_i:=\rho_i'.
$$

However, these operators do not satisfy the completeness condition:

$$
\rho := \sum_i \rho'_i \implies \text{tr}(\rho) = 1 \implies \rho \neq \mathbb{I}
$$

$$
M_i := \rho^{-\frac{1}{2}} \rho_i' \rho^{-\frac{1}{2}}
$$

### It is a projective (?) measurement which is defined as

$$
PGM = \{ |v_i\rangle \langle v_i| \}_i,
$$

where

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Our PGM is not necessarily projective!

### **Theorem (Barnum–Knill)**

$$
P^{pgm}(\mathcal{E}) \geq P^{opt}(\mathcal{E})^2.
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\sqrt{\text{Ppgm}(\mathcal{E})} \geq \text{P}^{\text{opt}}(\mathcal{E}).
$$

## **Gram matrix**

For a while, let's limit ourselves to the case where  $\rho_i$ 's are pure states. We can encode the inner product of all the states in an  $n \times n$  matrix *G*:

$$
G_{ij}=\sqrt{p_i p_j}\,\langle\psi_i\mid\psi_j\rangle
$$

$$
S := (\ket{\psi'_1}, \ldots, \ket{\psi'_n}) \implies G = S^{\dagger} S
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We may similarly encode the probability of getting outcome *i* and receiving state *j* in a matrix *P*:

$$
P_{i,j}:=\langle v_i|\psi_j'\rangle
$$

Then the success probability is

$$
P^{pgm}(\mathcal{E}) = \sum_{i=1}^{n} |\langle v_i | \psi'_i \rangle|^2 = \sum_{i=1}^{n} |P_{ii}|^2.
$$

## **Gram matrix and PGM**

We have:

$$
\left(P^2\right)_{ij} = \sum_{k=1}^n \left\langle \psi'_i \left| \rho^{-1/2} \right| \psi'_k \right\rangle \left\langle \psi'_k \left| \rho^{-1/2} \right| \psi'_j \right\rangle
$$

$$
= \left\langle \psi'_i \left| \left( \rho^{-1/2} \sum_{k=1}^n |\psi'_k\rangle \left\langle \psi'_k \right| \rho^{-1/2} \right) \right| \psi'_j \right\rangle
$$

$$
= G_{ij}
$$

Thus,

$$
P=\sqrt{G}.
$$

## **Corollary**

$$
P^{pgm}(\mathcal{E}) = \sum_{i=1}^{n} (\sqrt{G})_{ii}^{2}
$$

# <span id="page-14-0"></span>**[Two Lower Bounds for State](#page-14-0) [Discrimination](#page-14-0)**

In this part, we give the two lower bounds for the success probability of PGM:

- A bound obtained from the pairwise inner products
- A bound from the eigenvalues of the Gram matrix

#### **Lemma**

If for any 
$$
x > 0
$$
,  $\sqrt{x} \ge ax + bx^2$ , then  $(\sqrt{G})_{ii} \ge aG_{ii} + b\sum_{j=1}^n |G_{ij}|^2$ .

We find the parameters *a* and *b* such that  ${{aG_{ii}}} + b\sum\nolimits_{j = 1}^n {{{\left| {{G_{ij}}} \right|}^2}}$  is maximized.

The maximum is attained when  $a = \frac{3}{2\sqrt{r}}$  and  $b = -\frac{1}{2r^{3/2}}$ , where  $r =$  $\sum_{j=1}^n|\mathsf{G}_{ij}|^2$  $\frac{1+3i}{G_{ii}}$ .

Plugging it in our lemma:

$$
P^{pgm}(\mathcal{E}) \geq \sum_{i=1}^{n} \frac{p_i^2}{\sum_{j=1}^{n} p_j \left|\left\langle \psi_i \mid \psi_j \right\rangle\right|^2}
$$

## **A Bound from Eigenvalues**

$$
\sum_{i=1}^{n} (\sqrt{G})_{ii} = \sum_{i=1}^{n} \sqrt{\lambda_i}
$$

$$
\Rightarrow \left(\sum_{i=1}^{n} (\sqrt{G})_{ii}\right)^2 = \left(\sum_{i=1}^{n} \sqrt{\lambda_i}\right)^2
$$

$$
\Rightarrow n \sum_{i=1}^{n} (\sqrt{G})_{ii}^2 \ge \left(\sum_{i=1}^{n} \sqrt{\lambda_i}\right)^2
$$

$$
P^{pgm}(\mathcal{E}) \geq \frac{1}{n} \left( \sum_{i=1}^{n} \sqrt{\lambda_i} \right)^2
$$

Let  $\mathcal E$  be an ensemble of *n* mixed states  $\{\rho_i\}$  with a priori probabilities  $\{p_i\}$ , and having spectral decompositions  $\rho_i = \sum_{k=1}^d \lambda_{ik} |v_{ik}\rangle \langle v_{ik}|$ . Define F to be the ensemble of the *nd* pure states  $\{ |v_{ik}\rangle \}$  with a priori probabilities  $\{p_i \lambda_{ik}\}$ . Then  $P^{pgm}(\mathcal{E}) \geq P^{pgm}(\mathcal{F})$ .

# <span id="page-20-0"></span>**[Distinguishing Random Quantum](#page-20-0) [States](#page-20-0)**

## **Discrimination for Random Ensembles (Expectation)**

### **Theorem**

*Let E be an ensemble of n equiprobable d-dimensional quantum states*  $\{|\psi_i\rangle\}$  *with n*/*d*  $\rightarrow$  *r*  $\in$  (0*,*  $\infty$ ) *as n, d*  $\rightarrow \infty$ *, and let the components of |ψii in some basis be i.i.d. complex random variables with mean 0 and variance* 1*/d. Then*

$$
\mathbb{E}\left(P^{pgm}\left(\mathcal{E}\right)\right) \geq \begin{cases} \frac{1}{r}\left(1-\frac{1}{r}\left(1-\frac{64}{9\pi^2}\right)\right) & \text{if } n \geq d \\ 1-r\left(1-\frac{64}{9\pi^2}\right)^2 & \text{otherwise} \end{cases}
$$

*and in particular*  $\mathbb{E} \left( P^{pgm} (\mathcal{E}) \right) > 0.720$  when  $n \leq d$ .

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What is the point here?

- We have random ensembles.
- We need to bound the **expectation** of the probability of success.
- The premises of the theorem provide us the possibility of applying some nice results from random matrix theory.
- States are being chosen according to the Haar measure. 44

## **Discrimination for Random Ensembles (Concentration of Measure)**

#### **Theorem**

*Let E be an ensemble of n d-dimensional quantum states picked uniformly at random. Set*  $p = \mathbb{E}\left(P^{pgm}(\mathcal{E})\right) = \frac{1}{r}\left(1 - \frac{1}{r}\left(1 - \frac{64}{9\pi^2}\right)\right)$  *<i>if*  $n \geq d$ , and  $p = 1 - r\left(1 - \frac{64}{9\pi^2}\right)$  otherwise. Then

$$
\Pr\left[P^{pgm}(\mathcal{E}) \leq p - \epsilon\right] \leq 2 \exp\left(\frac{-C(2nd+1)\epsilon^2}{2}\right)
$$

*where*  $C = 1/(18\pi^3)$ .

# <span id="page-24-0"></span>**[Conclusion](#page-24-0)**

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- The importance of this work is:
	- finding analytic lower bounds for the success probability of pretty good measurements
	- using the theory of random matrices to apply the bounds in the case of random ensembles
- If my talk went well, you should probably know that:
	- what is the state discrimination problem.
	- the state discrimination problem has many applications.
	- pretty good measurement are indeed pretty good strategies.
	- obtaining lower (upper) bounds for the success probability of pretty good measurements might be useful to solve other problems.

<span id="page-26-0"></span>**[Thank you!](#page-26-0)**